

Lazy Functional Logic Programming and Encapsulated Search

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Not my own Work

Encapsulating Non-determinism in
Functional Logic Computations

Brassel, Hannus, Huch 2004

Computing with Subspaces

Antoy, Brassel 2007

Set Functions for Functional Logic Programming

Antoy, Hannus 2009

The Curry Language

- first-class nondeterminism
- call-by-need (lazy) semantics
- encapsulated search

First-class Non-determinism

coin :: Int
coin = 0 ? 1

cyi > coin

0

More? yes

1

More? yes

No more results

Example: Permutations

$\text{perm} :: [a] \rightarrow [a]$

$\text{perm } l =$

if null l then l

else insert (head l) (perm (tail l))

$\text{insert} :: a \rightarrow [a] \rightarrow [a]$

$\text{insert } x \text{ } xs = (x:xs) ?$

if null xs then fail

else head xs : insert x (tail xs)

Example : Permutations

`cyi > perm [1,2,3]`

`[1,2,3]`

More? *all*

`[1,3,2]`

`[2,1,3]`

`[2,3,1]`

`[3,1,2]`

`[3,2,1]`

First-class Nondeterminism

expressions can have
multiple values

interactive environment
for examining them

Infinitely Many Results

zeros :: Int

zeros = 0 ? zeros

$Cy_i > \text{zeros}$

0

More? yes

0

More? yes

0

More? no

Infinite Values

coins :: [Int]

coins = coin : coins

$c \gamma i \rangle$ coins

$\wedge C$

$c \gamma i \rangle$ head coins

0

More? all

1

Laziness

isSorted :: [Int] -> Bool

isSorted l =

if null l || null (tail l) then True

else head l <= head (tail l)

&& isSorted (tail l)

cyi> isSorted [0, -1..]

False

Lazy Nondeterminism

permsort :: [Int] → [Int]

permsort l =

let p = perm l in

if isSorted p then p else fail

same value

shared variable

Lazy Non-determinism

cyi > let $x = \text{coin}$ in $x + x$

0

More? yes

2

More? yes

No more results

$$0 + 0 = 0$$

~~$$0 + 1 = 1$$~~

~~$$1 + 0 = 1$$~~

$$1 + 1 = 2$$

Laziness

- infinitely many results
- infinite (intermediate) values
- evaluation order independence

let $x = a ? b$ in e



(let $x = a$ in e) ? (let $x = b$ in e)

Encapsulated Search

primitive operation

values :: $a \rightarrow [a]$

idea:

- values coin = $[0, 1]$
- head (values zeros) = σ
- values (a? b) \neq
values a ? values b

Encapsulated Search

1. Weak Encapsulation

2. Strong Encapsulation

3. Set Functions

Weak Encapsulation

may be nondeterministic

values coin = $[0, 1]$

but

let $x = \text{coin}$ in values $x = [0] ? [1]$

because x is introduced "outside"
of encapsulated expression

Weak Encapsulation

sharing between "inside" and "outside"

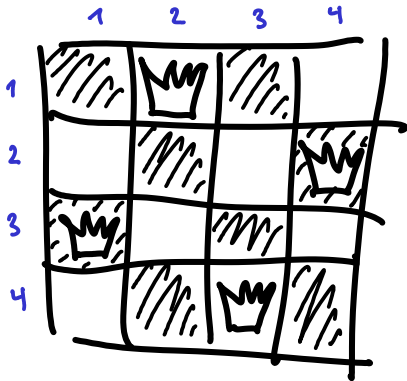
let x = coin in

values x ++ (x : values x)



$[0, 0, 0]$? $[1, 1, 1]$

N-Queens Problem



[3, 1, 4, 2]

16.5

N-Queens Problem

queens :: Int \rightarrow [Int]

queens n =

let p = perm [1..n] in

if null (values (capture p)) then p
else fail

capture :: [Int] \rightarrow ()

capture (_ ++ x : xs ++ y : _) =

if abs (x - y) == length xs + 1 then () else fail

Weak Encapsulation

different results based on
syntactic difference (scope)

not all choices encapsulated

Strong Encapsulation

encapsulates all choices

let $x = \text{coin}$ in values x



$[0, 1]$

although x is introduced "outside"
of encapsulated expression

Strong Encapsulation is Reusable

$\text{hasValue} :: a \rightarrow \text{Bool}$

$\text{hasValue } x = \text{not (null (values } x))$

$\text{firstValue} :: a \rightarrow a$

$\text{firstValue } x = \text{head (values } x)$

note: x is introduced "outside"

Strong Encapsulation

sharing between "inside" and "outside"

let $x = \text{coin}$ in values $x \text{ ++ } (x : \text{values } x)$

is not $[0, 1, 0, 0, 1] ? [0, 1, 1, 0, 1]$

but $[0, 1, 0, 0] ? [0, 1, 1, 1]$

(in all implementations
of strong encapsulation)

Strong Encapsulation

result depends on evaluation order

let $x = \text{coin}$ in values $x \# (x : \text{values } x)$

if x not yet evaluated

values $x = [0, 1]$

if x already evaluated to 0 (or 1)

values $x = [0]$ (or $[1]$)

N - Queens

queens :: Int → [Int]

queens n =

let p = perm [1..n] in

if p == p && null (values (capture p))

then p

else fail

↖ force evaluation
of p

Strong Encapsulation

- encapsulates all choices
- is reusable
- no evaluation-order independent implementation exists

Set Functions

No primitive values $:: a \rightarrow [a]$

Instead: set-valued variant
of every defined function

Set Functions

$\text{addCoin} :: \text{Int} \rightarrow \text{Int}$

$\text{addCoin } x = x + \text{coin}$

generates

$\text{addCoin}_{\text{set}} :: \text{Int} \rightarrow [\text{Int}]$ ¹

¹ conceptually: set, not list

Set Functions

encapsulate choices in body,
but not in arguments

$\text{addCoin}_{\text{Set}}(10 ? 20)$

$[10, 11] ? [20, 21]$

choice between 0 and 1

choice between 10 and 20

N - Queens

queens :: Int → [Int]

queens n =

let p = perm [1..n] in

if null (capture_{set} p) then p

else fail

Set Functions

similar to weak encapsulation

but separation of choices
based on argument-body
distinction
rather than on scoping

Lazy Functional Logic Programming

- first-class nondeterminism
- evaluation-order independent
call-by-need semantics
- interactive environment
for examining results
- encapsulated search

Weak Encapsulation

depends on scoping
not reusable

Strong Encapsulation

reusable

depends on evaluation order
(at least as implemented for Curry)

Set Functions

evaluation-order independent
no (reusable) strong encapsulation

Delimited Continuations ?

fail = []

$x ? y = \text{shift } k . k x ++ k y$

reset : weak or strong encapsulation ?

evaluation order independence ?

ありがとう