

A Haskell EDSL for (Parallel) Programming in Generate-Test-and-Aggregate Style

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Outline

1. Generate, Test, and Aggregate
2. Parallel Programming
3. A Haskell EDSL
4. Test-case generation?

Generate, Test, and Aggregate

- like generate-and-test
but with aggregate

data Item = Item { value :: Value, weight :: Weight }

knapsack :: Weight → [Item] → Value

knapsack maxW

= maximum . map (sum . map value)

· filter ((≤ maxW) . sum . map weight)

· sublists

Generator

generates a bag of lists

Test

discards some of those lists

Aggregator

reduces every list to single value

reduces bag of values to final result

Complexity

sublists $[1, 2, 3] =$

$\{ [], [1], [2], [3], [1, 2], [1, 3], [2, 3], [1, 2, 3] \}$

size of sublists l is $2^{\text{length } l}$

run time of knapsack max w l is

$O(2^{\text{length } l})$

Generators

only allowed to use certain operations
(polymorphic over some type class)

Test

by structural recursion on lists

Aggregators

by structural recursion on bag of lists

class Semiring S where

$zero, one :: S$

$(\oplus), (\otimes) :: S \rightarrow S \rightarrow S$

Laws

$$zero \oplus x = x \oplus zero = x \quad one \otimes x = x \otimes one = x$$

$$x \oplus (y \oplus z) = (x \oplus y) \oplus z \quad x \otimes (y \otimes z) = (x \otimes y) \otimes z$$

$$x \oplus y = y \oplus x$$

$$(x \oplus y) \otimes z = (x \otimes z) \oplus (y \otimes z)$$

$$x \otimes (y \oplus z) = (x \otimes y) \oplus (x \otimes z)$$

$$zero \otimes x = x \otimes zero = zero$$

instance Semiring $\mathcal{L}[a]$ where

$$\text{zero} = \mathcal{L} \emptyset$$

$$\text{one} = \mathcal{L} []$$

$$x \oplus y = x \uplus y \quad \text{-- bag union}$$

$$x \otimes y = \mathcal{L} [x_s ++ y_s \mid x_s \leftarrow x, y_s \leftarrow y]$$

Example

$$(\mathcal{L} [1] \oplus \mathcal{L} [2]) \otimes \mathcal{L} [3]$$

$$= \mathcal{L} [1], [2] \otimes \mathcal{L} [3] = \mathcal{L} [1,3], [2,3]$$

Another Instance

instance Semiring Value where

$$\text{zero} = -\infty$$

$$\text{one} = 0$$

$$x \oplus y = \max x y$$

$$x \otimes y = x + y$$

Type of Aggregators

maxValue :: [Item] → Value

Restricted Generator

$\text{sublists} :: [a] \rightarrow \mathcal{P}[a]$

$\text{sublists } \ell = \text{genSublists } (\lambda x \rightarrow \mathcal{P}[x]) \ell$

$\text{genSublists} :: \text{Semiring } s$

$\Rightarrow (a \rightarrow s) \rightarrow [a] \rightarrow s$

$\text{genSublists } _ [] = \text{one}$

$\text{genSublists } f (x:xs) =$

$(\text{one} \oplus f x) \otimes \text{genSublists } f xs$

Fusion

$$\begin{aligned} & \text{maxValue} . \text{genSublists } (\lambda x \rightarrow \llbracket x \rrbracket S) \\ = & \text{genSublists } (\lambda x \rightarrow \text{maxValue } \llbracket x \rrbracket S) \\ = & \text{genSublists } (\lambda x \rightarrow \text{value } x) \end{aligned}$$

Linear rather than exponential
run-time complexity

Generate, Test, and Aggregate

- intuitive, inefficient specification
- certain conditions allow fusion
(not shown: fuse aggregator and filter)
- efficient implementation
can be automatically derived

Parallel Programming

Divide and Conquer

Structural recursion
over monoid structure of lists

$$\text{sum } [] = \sigma$$

$$\text{sum } [x] = x$$

$$\text{sum } (xs ++ ys) = \text{sum } xs + \text{sum } ys$$

Parallel Sublist Generator

$\text{genSublists} :: \text{Semiring } s$
 $\Rightarrow (a \rightarrow s) \rightarrow [a] \rightarrow s$

$\text{genSublists } [] = \text{one}$

$\text{genSublists } f [x] = \text{one} \oplus f \ x$

$\text{genSublists } f (xs ++ ys) =$

$\text{genSublists } xs \otimes \text{genSublists } ys$

→ efficient, parallel knapsack function

How to obtain efficient implementation
from specification automatically?

GHC RULES pragma?

more reliable way?

Haskell EDSL

$\text{generate} :: (\forall s. \text{Semiring } s \Rightarrow (a \rightarrow s) \rightarrow [a] \rightarrow s)$
 $\rightarrow [a] \rightarrow \text{Gen } a$

$\text{test} :: \text{Monoid } m$
 $\Rightarrow (m \rightarrow \text{Bool}) \rightarrow (a \rightarrow m) \rightarrow \text{Gen } a \rightarrow \text{Gen } a$

$\text{aggregate} :: \text{Semiring } s$
 $\Rightarrow (a \rightarrow s) \rightarrow \text{Gen } a \rightarrow s$

Implementation

data Gen a where

Generate :: (∀s. Semiring s ⇒ (a → s) → [a] → s)
→ [a] → Gen a

Test :: Monoid m
⇒ (m → Bool) → (a → m) → Gen a → Gen a

generate = Generate
test = Test

Implementation of Fusion

aggregate :: Semiring s
 $\Rightarrow (a \rightarrow s) \rightarrow \text{Gen } a \rightarrow s$

aggregate f (Generate gen l) = gen f l

aggregate f (Test ok h gen) = ...

static guarantee:

no intermediate bags

Efficient knapsack function

knapsack :: Weight \rightarrow [Item] \rightarrow Value

knapsack maxW = aggregate value

- test ((\leq maxW).getSum) Sum
- generate genSublists

looks like exponential specification

Summary

- certain Gen-Test-Egg algorithms can be implemented efficiently
- parallelism orthogonal issue (depends on generator)
 - filters do not destroy parallelism
- implementation as Haskell EDSL strikingly simple

Test-case Generation ?

presented technique not restricted to lists
test can express (some) preconditions
aggregator can construct remaining test cases
complexity depends on number of results
not on number of discarded values (?)